

State Abstractions

What are abstractions and why study them?

- When we use more sophisticated function approximation, we are always generalizing the knowledge learned from one state to other similar states.
 - When is such generalization valid? **What states can be considered as “similar”?**
- To answer these questions, it is worth studying the simplest form of generalization: abstractions
- State abstractions \approx aggregate equivalent (or similar) states and run tabular algorithms

Examples of state abstractions

- Multiple ways of expressing an abstraction
 - Mapping ϕ from original (or *raw*) states to abstract states
 - Partition over the state space
 - An equivalence notion over raw states
- Example 1: discretize a continuous state space
 - Mapping from continuous state to the grid
 - Partition is obvious
 - Two original states are equivalent if they fall in the same grid
- Example 2: Suppose the original state is described by some state variables $s = (x, y)$. $\phi(s) = x$ is an abstraction
 - mapping $\phi : (x, y) \mapsto x$
 - Partition over $\{(x, y)\}$
 - $s_1 = (x_1, y_1)$ is equiv to $s_2 = (x_2, y_2)$ iff $x_1 = x_2$ (i.e., $\phi(s_1) = \phi(s_2)$)

Notations and Formal Setup

- MDP $M = (S, A, P, R, \gamma)$
- Abstraction $\phi : S \rightarrow S_\phi$
- Surjection — aggregate states and treat as equivalent
- Are the aggregated states really equivalent?
- Do they have the same...
 - optimal action?
 - Q^* values?
 - dynamics and rewards?

Abstraction hierarchy

An abstraction ϕ is ... if ... $\forall s^{(1)}, s^{(2)}$ where $\phi(s^{(1)}) = \phi(s^{(2)})$

- π^* -irrelevant: $\exists \pi_M^*$ s.t. $\pi_M^*(s^{(1)}) = \pi_M^*(s^{(2)})$
- Q^* -irrelevant: $\forall a, Q_M^*(s^{(1)}, a) = Q_M^*(s^{(2)}, a)$
- Model-irrelevant: $\forall a \in A,$
 (bisimulation) $\forall a \in A, x' \in S_\phi,$

$$R(s^{(1)}, a) = R(s^{(2)}, a)$$

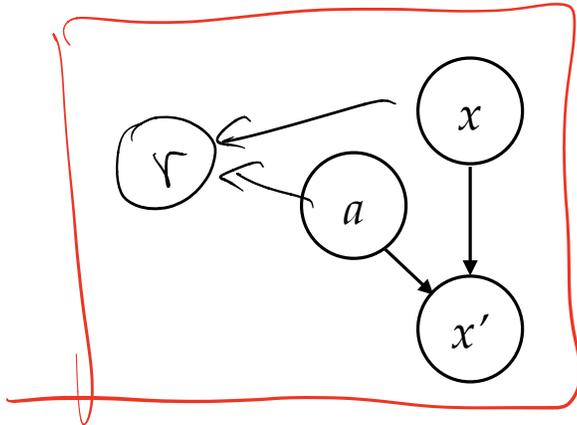
$$P(x' | s^{(1)}, a) = P(x' | s^{(2)}, a)$$



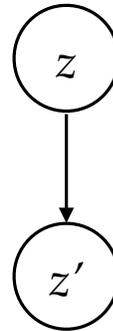
$$\sum_{s' \in \phi^{-1}(x')} P(s' | s^{(1)}, a)$$

Theorem: Model-irrelevance $\Rightarrow Q^*$ -irrelevance $\Rightarrow \pi^*$ -irrelevance

Why not $P(s' \mid s^{(1)}, a) = P(s' \mid s^{(2)}, a)$?



MDP M



Markov chain C

$$\phi: (x, z) \mapsto x.$$

$(x, z^{(1)})$ and $(x, z^{(2)})$ cannot be aggregated under the s' -based condition

$$P((x', z') \mid (x, z), a) = P_M(x' \mid x, a) \cdot P_C(z' \mid z)$$

integrated out by bisimulation

"Overly strict criteria" note $s = (x, z)$ $s' = (x', z')$.

$$\frac{(x, z^{(1)})}{s^{(1)}}, \frac{(x, z^{(2)})}{s^{(2)}}$$

$$P(s' | s^{(1)}, a) = P_M(x' | x, a) \cdot P_C(z' | z^{(1)})$$

$$P(s' | s^{(2)}, a) = P_M(x' | x, a) \cdot P_C(z' | z^{(2)})$$

identical.

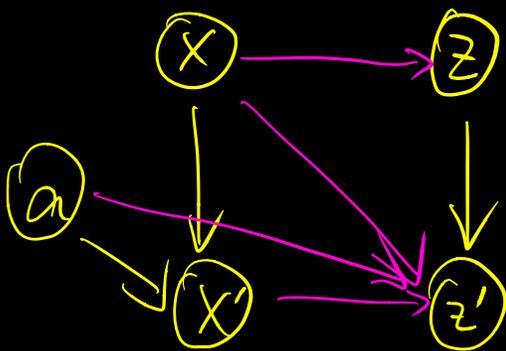
Bisimulation reward criterion ✓.

$$\sum_{s' \in \phi^{-1}(x')} P(s' | s^{(1)}, a) = \sum_{z'} P_M(x' | x, a) \cdot P_C(z' | z^{(1)})$$

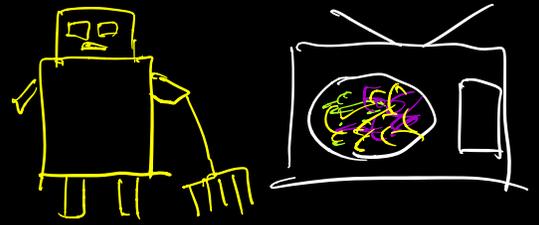
$$= P_M(x' | x, a) \cdot \sum_{z'} P_C(z' | z^{(1)}) = P_M(x' | x, a)$$

$$\sum_{s' \in \phi^{-1}(x')} P(s' | s^{(2)}, a) = \underline{\underline{P_M(x' | x, a)}}$$

Optional: add maximum arrows so that ϕ is still bisimulation.



"Noisy TV" Problem.



Naive model-based RL:

predict (x', z') given $(x, z), a$.

state = (x, z)

$R((x, z), a)$ only depends on x , and a .

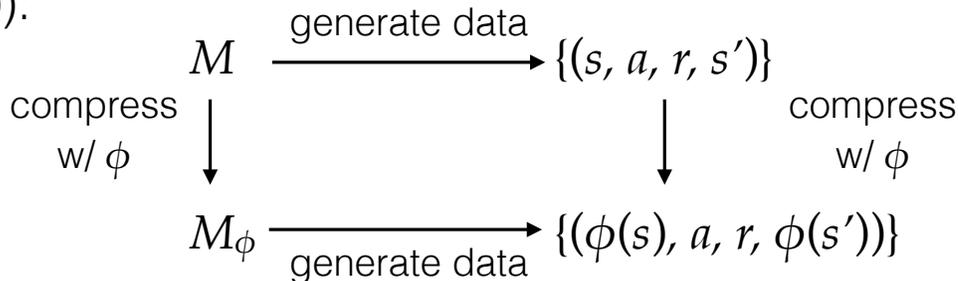
Robot & rest of the home.

TV content

The abstract MDP implied by bisimulation

ϕ is bisimulation: $R(s^{(1)}, a) = R(s^{(2)}, a)$, $P(x' \mid s^{(1)}, a) = P(x' \mid s^{(2)}, a)$

- MDP $M_\phi = (S_\phi, A, P_\phi, R_\phi, \gamma)$
- For any $x \in S_\phi, a \in A, x' \in S_\phi$
 - $R_\phi(x, a) = R(s, a)$ for any $s \in \phi^{-1}(x)$
 - $P_\phi(x' \mid x, a) = P(x' \mid s, a)$ for any $s \in \phi^{-1}(x)$
- No way to distinguish between the two routes (if a only depends on $\phi(s)$):



Bisimulation $\Rightarrow Q^*$ -irrelevance

$$[Q_{M_\phi}^*]_M : (s, a) \mapsto Q_{M_\phi}^*(\phi(s), a)$$

- Consider the Q^* in M_ϕ , $Q_{M_\phi}^*$ (dimension: $|S_\phi \times A|$)
- *Lift* $Q_{M_\phi}^*$ from S_ϕ to S (populate aggregated states with the same value)
- Useful notation: Φ is a $|S_\phi| \times |S|$ matrix, with
$$\Phi(x, s) = \mathbb{1}[\phi(s) = x]$$
 - lifting a state-value function: $[V_{M_\phi}^*]_M = \Phi^\top V_{M_\phi}^*$
 - collapsing the transition distribution: $\Phi P(s, a)$
- Claim: $[Q_{M_\phi}^*]_M = Q_M^*$ (proof)

$$\Phi_{\times \phi} = \begin{matrix} |S| \\ \downarrow \\ |S| \times |S| \end{matrix} \begin{matrix} |S| \\ \downarrow \\ |S| \times |S| \end{matrix} \begin{matrix} |S| \\ \downarrow \\ |S| \times |S| \end{matrix} = \begin{matrix} |S| \\ \downarrow \\ |S| \times |S| \end{matrix}$$

Bisimulation: $\forall s^{(1)}, s^{(2)}, \text{ s.t. } \phi(s^{(1)}) = \phi(s^{(2)})$

• $\forall a, R(s^{(1)}, a) = R(s^{(2)}, a)$

• $\forall a, \underbrace{\Phi}_{|S_\phi| \times |S|} \underbrace{P(\cdot | s, a)}_{|S| \times |S|} = \underbrace{\Phi}_{|S_\phi| \times |S|} \underbrace{P(\cdot | s, a)}_{|S| \times |S|}$

Proof: $\underbrace{[Q_{M_\phi}^*]_M}_{\downarrow} = Q_M^* \iff \underbrace{[Q_{M_\phi}^*]_M}_{\downarrow} = \mathcal{J}_M [Q_{M_\phi}^*]_M$

$\iff \forall (s, a), \text{ RHS}(s, a)$

$= R(s, a) + \gamma \mathbb{E}_{s' \sim P(\cdot | s, a)} [[V_{M_\phi}^*]_M(s')]$

$= R(s, a) + \gamma \langle P(\cdot | s, a), \Phi^T V_{M_\phi}^* \rangle$

$= R_\phi(\phi(s), a) + \gamma \langle \underbrace{\Phi P(\cdot | s, a)}_{\downarrow}, V_{M_\phi}^* \rangle$

Reason: $\langle x, A^T y \rangle = \underline{x^T A^T y} = (Ax)^T y = \langle Ax, y \rangle$

$= R_\phi(\phi(s), a) + \gamma \mathbb{E}_{x' \sim P_\phi(\phi(s), a)} [V_{M_\phi}^*(x')]$

$= Q_{M_\phi}^*(\phi(s), a) = \text{LHS}(s, a) \quad \text{Q.E.D. } \square$

$(\mathcal{J}f)(s, a) = R(s, a) + \gamma \mathbb{E}_{s' \sim P(\cdot | s, a)} [\max_a f(s, a)]$

$\Phi^T x = \begin{bmatrix} \phi(s) \\ \phi(s) \\ \phi(s) \\ \vdots \end{bmatrix} \times \begin{bmatrix} x \\ x \\ x \\ \vdots \end{bmatrix}$

Useful/fun facts about bisimulation

- Q_M^π is preserved *for any π lifted from an abstract policy*
- Given any lifted π , distribution over reward sequence is preserved (assuming reward is deterministic function of s, a)
- Coarsest bisimulation always exists: in any MDP, the common coarsening of two bisimulations is always a bisimulation
 - e.g., ϕ_1 tells you to ignore some state variables, ϕ_2 tells you to ignore some others \Rightarrow can ignore both sets of variables!
 - Intuitive but nontrivial; needs proof (see notes)

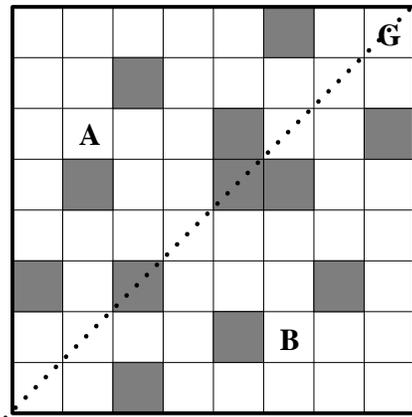
Useful/fun facts about bisimulation

- Recall that bisimulation is defined by a reward condition and a transition condition
- Guess what's the coarsest bisimulation if we drop the reward condition and only require the transition condition?
 - Aggregate all states together!
 - reward function defines a notion of (short-term) saliency
 - can extend the definition by replacing reward function with other functions (even not real-valued ones) whose codomain is equipped with an equivalence notion

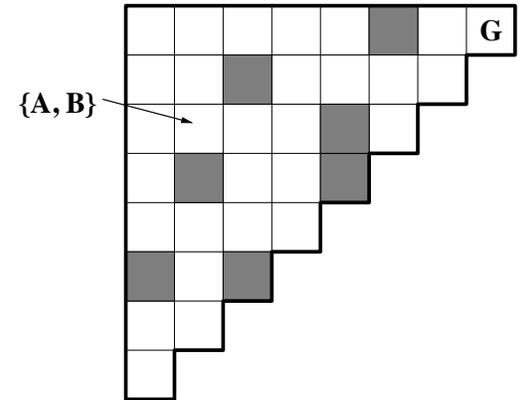
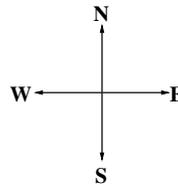
The abstract model

- Consider planning, e.g., want to plan in the abstract model instead of the original model to reduce computation cost
- Approach: compress the model (M_ϕ), and plan in M_ϕ (and lift the policy back to M)
- We already showed: if ϕ is bisimulation, this approach produces an optimal policy of M
- What if ϕ is Q^* -irrelevant? or π^* -irrelevant?
- π^* -irrelevant: learned policy can be suboptimal (see refs in Li et al'06)
- Q^* -irrelevant: surprisingly, optimality is preserved; for details and further reading, see ref notes.

Extension to handle action aggregation/permutation: Homomorphisms



(a)



(b)

Figure from: Ravindran & Barto. Approximate Homomorphisms: A framework for non-exact minimization in Markov Decision Processes. 2004.

Approximate abstractions

1. ϕ is an ϵ_{π^*} -approximate π^* -irrelevant abstraction, if there exists an abstract policy $\pi : \phi(\mathcal{S}) \rightarrow \mathcal{A}$, such that $\|V_M^* - V_M^{[\pi]} \|_\infty \leq \epsilon_{\pi^*}$.
2. ϕ is an ϵ_{Q^*} -approximate Q^* -irrelevant abstraction if there exists an abstract Q -value function $f : \phi(\mathcal{S}) \times \mathcal{A} \rightarrow \mathbb{R}$, such that $\|[f]_M - Q_M^* \|_\infty \leq \epsilon_{Q^*}$.
3. ϕ is an (ϵ_R, ϵ_P) -approximate model-irrelevant abstraction if for any $s^{(1)}$ and $s^{(2)}$ where $\phi(s^{(1)}) = \phi(s^{(2)})$, $\forall a \in \mathcal{A}$,

$$|R(s^{(1)}, a) - R(s^{(2)}, a)| \leq \epsilon_R, \quad \left\| \Phi P(s^{(1)}, a) - \Phi P(s^{(2)}, a) \right\|_1 \leq \epsilon_P. \quad (3)$$

Theorem 2. (1) If ϕ is an (ϵ_R, ϵ_P) -approximate model-irrelevant abstraction, then ϕ is also an approximate Q^* -irrelevant abstraction with approximation error $\epsilon_{Q^*} = \frac{\epsilon_R}{1-\gamma} + \frac{\gamma\epsilon_P R_{\max}}{2(1-\gamma)^2}$.

(2) If ϕ is an ϵ_{Q^*} -approximate Q^* -irrelevant abstraction, then ϕ is also an approximate π^* -irrelevant abstraction with approximation error $\epsilon_{\pi^*} = 2\epsilon_{Q^*}/(1-\gamma)$.